

# A Theory of Second-Order Wireless Network Optimization and Its Application on Aol

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# Challenges with Current Methods

- Network Utility Maximization (NUM) showed success in finding optimal solutions for traditional objective functions.
- Current applications for real-time sensory estimation and video streaming use other metrics such as Age of Information (AoI).
- NUM techniques almost always fail when optimizing for newer performance metrics such as Age of Information (AoI) and Quality of Experience (QoE).
- NUM techniques such as primal-dual decomposition or Lyapunov drift-plus penalty **only capture** first order-statistics (i.e. mean).

# Our Approach: Second-Order Optimization



- We can capture short-term system behavior by observing both the mean and temporal variance of random delivery processes.
  - Newer metrics involve higher-order behaviors to capture short-term network behavior.
- Given channel models, we derive the second-order capacity region consisting of all mean and temporal variance of packet delivery processes.
- We propose a new scheduling policy called Variance-Weighted Deficit (VWD), and prove it achieves every inner-point in the second-order capacity region.
- We apply VWD on an open problem: Optimizing AoI over independent Gilbert-Elliott channels.

# Wireless System Framework

- One Access Point (AP) serving  $N$  clients referenced as  $i = 1, 2, \dots, N$  at timeslots  $t = 1, 2, \dots$
- Each client is associated with an ON-OFF channel. The AP observes the channels at each timestep.
- AP can choose one client with an ON-channel for transmission at time  $t$ .

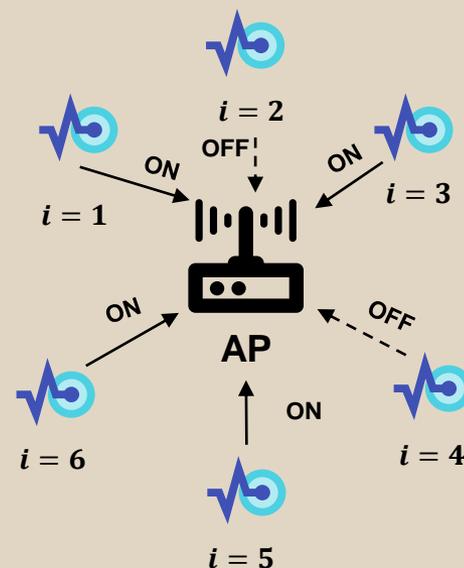


Fig. 1: AP serving six clients.

# Second-Order Model of Wireless Channels



- Given a subset of clients  $S \subseteq \{1, 2, \dots, N\}$ . Let  $X_S(t)$  be the indicator function that at least one client has an ON channel at time  $t$ .

- Denote the channel mean rate of  $S$  as

$$m_S := \lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T X_S(t)}{T}.$$

- Denote the channel temporal variance of  $S$  as

$$v_S^2 := E \left[ \left( \lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T X_S(t) - T m_S}{\sqrt{T}} \right)^2 \right].$$

- Define the second-order channel model as a collection of all the mean and temporal variance of subsets

$$\{(m_S, v_S^2) \mid S \subseteq \{1, 2, \dots, N\}\}.$$

- How to define a client's mean rate and temporal variance?

# Second-Order Model for Application Performance



- Let  $Z_i(t)$  be the indicator function that client  $i$  receives a packet at time  $t$ .

- Define client  $i$  service mean rate as  $\mu_i := \lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T Z_i(t)}{T}$ .

- Define client  $i$  service temporal variance as

$$\sigma_i^2 := E \left[ \left( \lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T Z_i(t) - T \mu_i}{\sqrt{T}} \right)^2 \right].$$

- Utility of client  $i$  is a function of  $(\mu_i, \sigma_i^2)$ , denoted by  $F_i(\mu_i, \sigma_i^2)$ .
- Goal to maximize  $\sum_{i=1}^N F_i(\mu_i, \sigma_i^2)$ , given network constraints.
  - How to define the network constraints?

# Second-Order Model of Network Constraints



- Given a second-order channel model  $\{(m_s, v_s^2) \mid S \subseteq \{1, 2, \dots, N\}\}$  at time  $t$ .
- Define the second-order capacity region as the set of all  $\{(\mu_i, \sigma_i^2) \mid 1 \leq i \leq N\}$  such that a policy achieves
  - $\lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T Z_i(t)}{T} = \mu_i$ .
  - $E\left[\left(\lim_{T \rightarrow \infty} \frac{\sum_{t=1}^T Z_i(t) - T \mu_i}{\sqrt{T}}\right)^2\right] \leq \sigma_i^2$  for all clients  $i = 1, 2, \dots, N$ .

# Outer-Bound of Second-Order Capacity Region



- **Theorem.** The second-order delivery model **can be in** the capacity region if the following holds
- Sum of clients' mean rate in the subset  $S$  is **less than or equal to** channel mean  $\sum_{i \in S} \mu_i \leq m_S$  for all  $S \subseteq \{1, 2, \dots, N\}$ .
- Sum of all clients' mean rate sum to the channel mean rate  $\sum_{i=1}^N \mu_i = m_{\{1, 2, \dots, N\}}$ .
- Sum of clients' temporal variance square root is bigger or equal to channel temporal variance  $\sum_{i=1}^N \sqrt{\sigma_i^2} \geq \sqrt{v_{\{1, 2, \dots, N\}}^2}$ .
- Client mean rate  $\mu_i \geq 0$  for all  $i$ .

# Inner-Bound of Second-Order Capacity Region



- **Theorem.** The second-order delivery model **is in** the capacity region if the following holds
- Sum of clients' mean rate in the subset  $S$  **is less than** channel mean rate  $\sum_{i \in S} \mu_i < m_S$  for all  $S \subsetneq \{1, 2, \dots, N\}$ .
- Sum of all clients' mean rate sum to the channel mean rate  $\sum_{i=1}^N \mu_i = m_{\{1, 2, \dots, N\}}$ .
- Sum of clients' temporal variance is bigger or equal to the channel temporal variance  $\sum_{i=1}^N \sqrt{\sigma_i^2} \geq \sqrt{v_{\{1, 2, \dots, N\}}^2}$ .
- Client mean rate  $\mu_i \geq 0$  and client temporal variance  $\sigma_i^2 > 0$  for all  $i$ .
  - How to achieve this inner-bound?

# Variance-Weighted Deficit (VWD) Policy



- **Theorem.** VWD policy achieves every point in the inner bound of second-order capacity region.
- Given a point  $(\mu_i, \sigma_i^2)$  in the bound.
- At time  $t$ , define client  $i$  deficit as  $d_i(t) = t\mu_i - \sum_{\tau=1}^t Z_i(\tau)$ .
- From clients with ON channels, the controller picks the client with the largest  $\frac{d_i(t-1)}{\sqrt{\sigma_i^2}}$ .
- Since we are interested in Aol performance, how can we measure it for a policy such as VWD?

# Problem we Consider

- Optimize Aol by deriving the second-order model over Gilbert-Elliot channels.

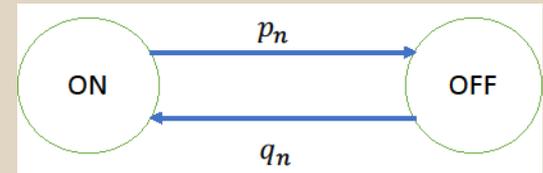


Fig. 2: Gilbert-Elliot channel.

- **Gilbert-Elliot channel:** two-state Markov process with transition probabilities  $p_i$  and  $q_i$ .
- Each client generates an update with probability  $\lambda_i$  and only keeps the newest update in memory.
- **Age of Information (Aol):** time difference between the newest information update at the source and the delivered information at the destination.

# Second-Order Model for Gilbert-Elliott Channels



- Under the Gilbert-Elliott channels, for all  $S$

$$m_S = 1 - \prod_{i \in S} \frac{p_i}{p_i + q_i},$$

$$v_S^2 = 2 \sum_{k=1}^{\infty} \left( \prod_{i \in S} G_i(k+1) \prod_{i \in S} \frac{p_i}{p_i + q_i} \right) \prod_{i \in S} \frac{p_i}{p_i + q_i} + \prod_{i \in S} \frac{p_i}{p_i + q_i} - \left( \prod_{i \in S} \frac{p_i}{p_i + q_i} \right)^2,$$

$$\text{with } G_i(k) = \frac{p_i}{p_i + q_i} + \frac{q_i}{p_i + q_i} (1 - p_i - q_i)^{k-1}.$$

- Optimize Aol over Gilbert-Elliott channel model. How to express Aol using the second-order model?

# Second-Order Expression of Aol of $(\mu_i, \sigma_i^2)$



- Let  $B_i(n)$  be the time between  $n^{th}$  and  $n + 1^{th}$  deliveries.
- Long-term average Aol (theoretical Aol)  $\overline{Aol}_i$  is given as
$$\overline{Aol}_i = \frac{E[B_i^2]}{2E[B_i]} + \frac{1}{\lambda_i} - \frac{1}{2}.$$
- Estimate  $\overline{Aol}_i$  from a Brownian motion random process  $BM_{\mu_i, \sigma_i^2}(t)$ .
- Approximate  $B_i(n)$  by the amount of time the Brownian process increases by 1.

# Second-Order Expression of Aol of $(\mu_i, \sigma_i^2)$



- Therefore, we can approximate  $\overline{Aol}_i$  (empirical Aol) as

$$\overline{Aol}_i \approx \frac{1}{2} \left( \frac{\sigma_i^2}{\mu_i^2} + \frac{1}{\mu_i} \right) + \frac{1}{\lambda_i} - \frac{1}{2}.$$

- Second-order optimization problem involves finding policy that maximizes  $\sum_{i=1}^N F_i(\mu_i, \sigma_i^2)$ .
  - For Gilbert-Elliott channels, what is the Aol performance function?

# Finding Clients' $(\mu_i, \sigma_i^2)$ for VWD

- With the goal of minimizing AoI over Gilbert-Elliot channels, we define objective function for client  $i$  as

$$F_i(\mu_i, \sigma_i^2) = -\frac{1}{2} \left( \frac{\sigma_i^2}{\mu_i^2} + \frac{1}{\mu_i} \right) - \frac{1}{\lambda_i} + \frac{1}{2}.$$

- We obtain the optimal delivery model for our VWD policy using steps:
  - Find all sets of the second-order channel model  $\{(m_s, v_s^2) \mid S \subsetneq \{1, 2, \dots, N\}\}$ .
  - Calculate client mean rates that satisfy  $\sum_{i \in S} \mu_i \leq m_s - \delta$  and  $\sum_{i=1}^N \mu_i = m_{\{1, 2, \dots, N\}}$ .
  - Client's temporal variance is lower bounded by channel temporal variance  $\sum_{i=1}^N \sqrt{\sigma_i^2} \geq \sqrt{v_{\{1, 2, \dots, N\}}^2}$ .
  - $\mu_i \geq 0$  and  $\sigma_i^2 > 0$  for all  $i$ .

# AoI Estimation of a Single Client

- Evaluate the theoretical AoI and empirical AoI on a single client.
- Results averaged over a 1000 independent runs.
- Each run contains 50000 timeslots.
- Empirical AoI is almost identical to the theoretical AoI.

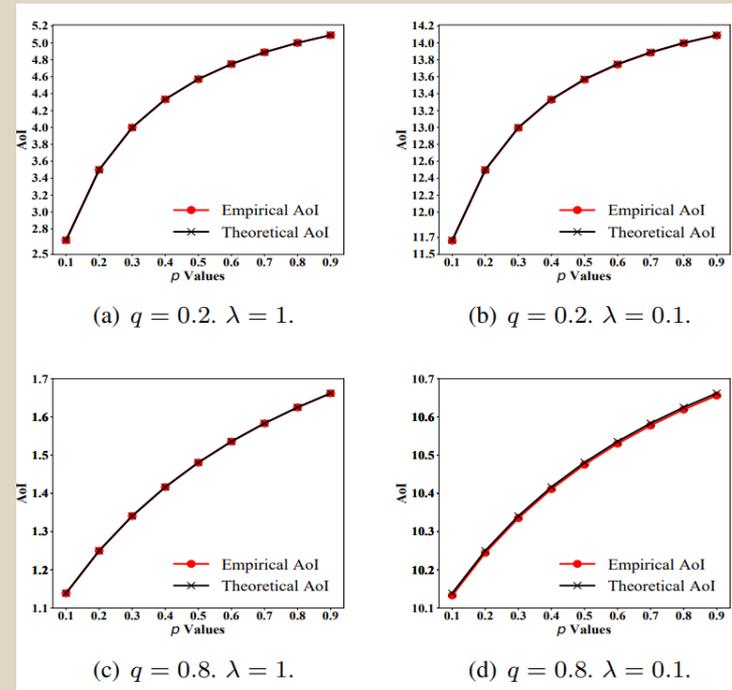
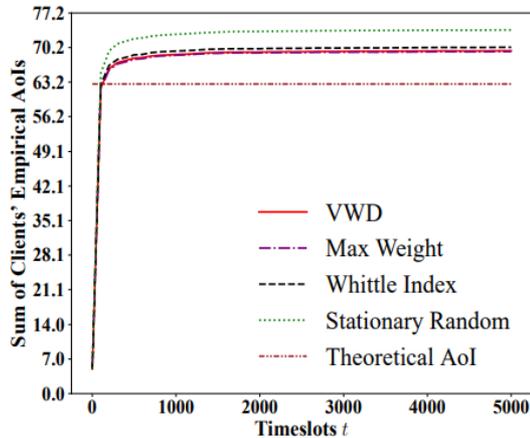


Fig. 3: Model validation for a single client.

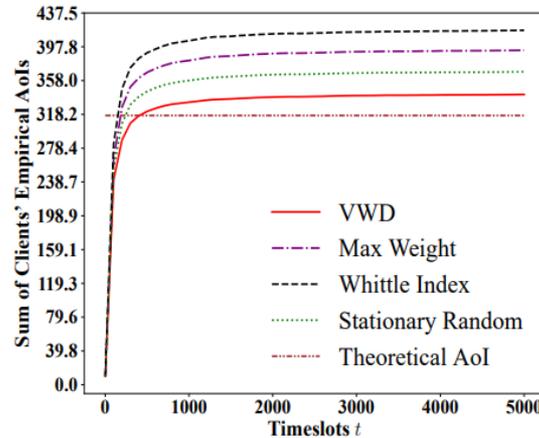
# Simulations Setting

- Compare VWD against policies:
  - **Whittle index:** schedules the highest-indexed ON client
$$W_i(t) = \frac{AoI_i^2(t)}{2} - \frac{AoI_i(t)}{2} + \frac{AoI_i(t)}{q_i/(p_i+q_i)}.$$
  - **Stationary randomized:** picks an ON client randomly proportional to  $\mu_i$ .
  - **Max weight:** picks an ON client with the largest  $\frac{AoI_i(t) - z_i(t)}{\mu_i}$  with
$$z_i(t) = \frac{1}{\lambda_i}.$$
- Simulations are ran for 1000 independent runs for 5000 timeslots.
- $\lambda_i$  randomly chosen from the range  $(\frac{0.1}{N}, \frac{1}{N})$ .
- Objective is to minimize  $\sum_i \alpha_i \overline{AoI_i}$ , with client weight  $\alpha_i$ .

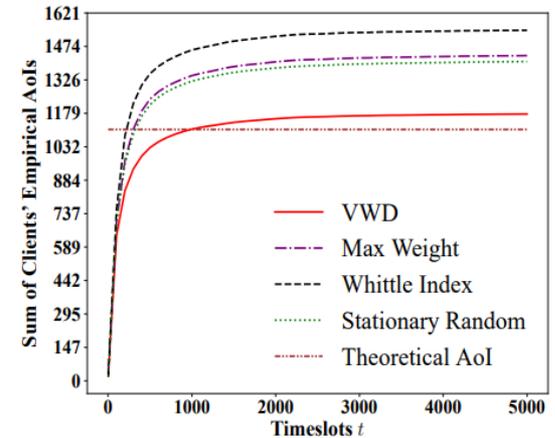
# Aol with Equal Weights' Results



(a)  $N = 5$  Clients.



(b)  $N = 10$  Clients.

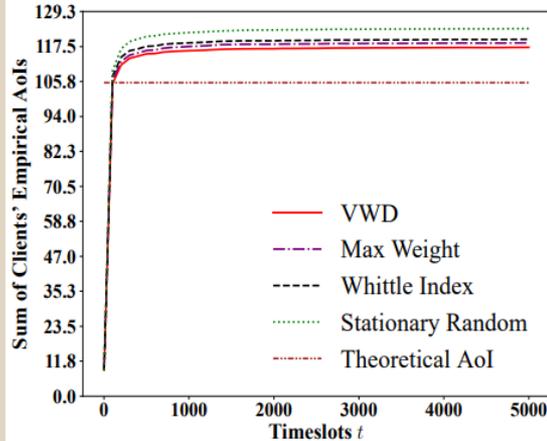


(c)  $N = 20$  Clients.

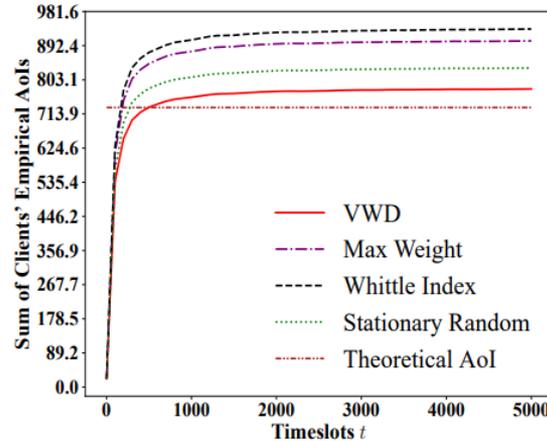
Fig. 4: Uniform empirical Aol results averaged over 1000 runs.

- Three different systems with: 5, 10, and 20 clients.
- VWD outperforms other policies for  $N = 10$  and  $N = 20$ . VWD performs similar to Max weight for  $N = 5$ .
- VWD's empirical Aol is close to the theoretical Aol compared to other policies. VWD Aol approximation is accurate.

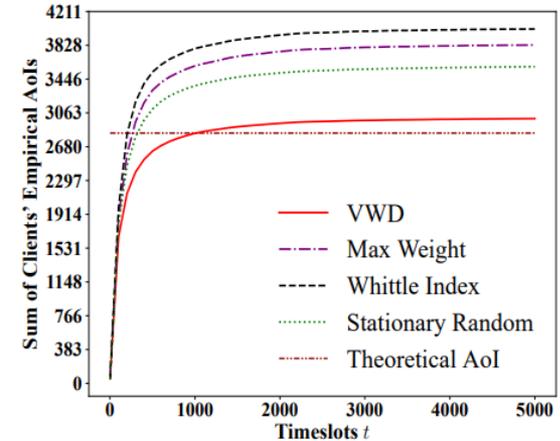
# Weighted Aol Results



(a)  $N = 5$  Clients.



(b)  $N = 10$  Clients.



(c)  $N = 20$  Clients.

Fig. 5: Weighted empirical Aol results averaged over 1000 runs.

- Weights  $\alpha_i$  were randomly selected from the range (1,5).
- VWD outperforms other scheduling policies in the weighted Aol setting.
- VWD's empirical Aol is close to the theoretical Aol compared to other policies. VWD Aol approximation is accurate.

# Summary

- Proposed a new general model: the second-order capacity region for wireless networks.
- Introduced a new scheduling policy, VWD, that captures second-order statistics (temporal variance) within a second-order capacity region.
- Applied VWD on the unsolved optimization problem over Gilbert-Elliott channels.
- VWD outperforms other compared scheduling policies in both the weighted and unweighted Aol settings.

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